## **AMENDMENTS TO THE CLAIMS**

This listing of claims will replace all prior versions, and listings, of claims in the application:

## **Listing of Claims:**

| 1  | 1. (Currently amended) A method for verifying type safety of an                  |
|----|--|
| 2  | application snapshot, the application snapshot including a state of an executing |
| 3  | program that is moved from a first computing device to a second computing        |
| 4  | device across a network in order to continue execution on the second computing   |
| 5  | device, the method comprising:   |
| 6  | receiving the application snapshot of the executing program from the first       |
| 7  | computing device on the second computing device, wherein the application         |
| 8  | snapshot includes a subprogram, an operand stack, and a point of execution;      |
| 9  | examining the application snapshot on the second computing device to             |
| 10 | identify the subprogram being executed and the point of execution within the     |
| 11 | subprogram;  |
| 12 | examining the subprogram on the second computing device to determine             |
| 13 | an expected structure of the operand stack at the point of execution;            |
| 14 | validating that the state of the application snapshot on the second              |
| 15 | computing device is consistent with the expected structure of the operand stack; |
| 16 | verifying on the second computing device that variables and argument             |
| 17 | arguments within the application snapshot are of the proper type; and            |
| 18 | if the state of the application snapshot is validated as consistent with the     |
| 19 | expected structure of the operand stack, resuming execution of the application   |
| 20 | snapshot on the second computing device at the point of execution on the first   |
| 21 | computing device.  |

| 1 | 2. (Original) The method of claim 1, wherein examining the subprogram                 |
|---|---|
| 2 | to determine the expected structure of the operand stack at the point of execution    |
| 3 | involves examining the subprogram with a code verifier, wherein the code verifier     |
| 4 | ensures that:   |
| 5 | the subprogram does not cause the operand stack to overflow and                       |
| 6 | underflow;  |
| 7 | a use of a local variable does not violate type safety; and                           |
| 8 | an argument of an instruction is of an expected type.                                 |
| 1 | 3. (Original) The method of claim 1, wherein the operand stack contains at            |
| 2 | least one local variable, at least one argument that is passed as a parameter to the  |
| 3 | subprogram, and an offset to the point of execution within the subprogram.            |
| 1 | 4. (Original) The method of claim 2, wherein the expected structure of the            |
| 2 | operand stack includes a collective size of entries and the types of entries expected |
| 3 | on the operand stack at the point of execution within the subprogram.                 |
| 1 | 5 (Canceled).   |
| 1 | 6. (Original) The method of claim 4, wherein validating that the state of             |
| 2 | the application snapshot on the second computing device is consistent with the        |
| 3 | expected structure of the operand stack involves ensuring that the collective size    |
| 4 | of entries and the types of entries on the operand stack agree with the collective    |
| 5 | size of entries and the types of entries expected on the operand stack.               |
| 1 | 7. (Original) The method of claim 1, wherein resuming execution of the                |
| 2 | application snapshot involves restarting the subprogram at the point of execution     |

within the second computing device.

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| 1  | 8. (Currently amended) A computer-readable storage medium storing                |
|----|--|
| 2  | instructions that when executed by a computer causes the computer to perform a   |
| 3  | method for verifying type safety of an application snapshot, the application     |
| 4  | snapshot including a state of an executing program that is moved from a first    |
| 5  | computing device to a second computing device across a network in order to       |
| 6  | continue execution on the second computing device, the method comprising:        |
| 7  | receiving the application snapshot of the executing program from the first       |
| 8  | computing device on the second computing device, wherein the application         |
| 9  | snapshot includes a subprogram, an operand stack, and a point of execution;      |
| 10 | examining the application snapshot on the second computing device to             |
| 11 | identify the subprogram being executed and the point of execution within the     |
| 12 | subprogram;  |
| 13 | examining the subprogram on the second computing device to determine             |
| 14 | an expected structure of the operand stack at the point of execution;            |
| 15 | validating that the state of the application snapshot on the second              |
| 16 | computing device is consistent with the expected structure of the operand stack; |
| 17 | verifying on the second computing device that variables and argument             |
| 18 | arguments within the application snapshot are of the proper type; and            |
| 19 | if the state of the application snapshot is validated as consistent with the     |
| 20 | expected structure of the operand stack, resuming execution of the application   |
| 21 | snapshot on the second computing device at the point of execution from the first |
| 22 | computing device.  |

9. (Original) The computer-readable storage medium of claim 8, wherein examining the subprogram to determine the expected structure of the operand stack at the point of execution involves examining the subprogram with a code verifier, wherein the code verifier ensures that:

| 5 | the subprogram does not cause the operand stack to overflow and                       |
|---|---|
| 6 | underflow;  |
| 7 | a use of a local variable does not violate type safety; and                           |
| 8 | an argument of an instruction is of an expected type.                                 |
| 1 | 10. (Original) The computer-readable storage medium of claim 8, wherein               |
| 2 | the operand stack contains at least one local variable, at least one argument that is |
| 3 | passed as a parameter to the subprogram, and an offset to the point of execution      |
| 4 | within the subprogram.  |
| 1 | 11. (Original) The computer-readable storage medium of claim 9, wherein               |
| 2 | the expected structure of the operand stack includes a collective size of entries and |
| 3 | the types of entries expected on the operand stack at the point of execution within   |
| 4 | the subprogram.   |
| 1 | 12 (Canceled).  |
| 1 | 13. (Original) The computer-readable storage medium of claim 11,                      |
| 2 | wherein validating that the state of the application snapshot on the second           |
| 3 | computing device is consistent with the expected structure of the operand stack       |
| 4 | involves ensuring that the collective size of entries and the types of entries on the |
| 5 | operand stack agree with the collective size of entries and the types of entries      |
| 6 | expected on the operand stack.  |
| 1 | 14. (Original) The computer-readable storage medium of claim 8, wherein               |
| 2 | resuming execution of the application snapshot involves restarting the subprogram     |
| 3 | at the point of execution within the second computing device.                         |

| 1  | 15. (Currently amended) An apparatus that facilitates verifying type safety           |
|----|---|
| 2  | of an application snapshot, the application snapshot including a state of an          |
| 3  | executing program that is moved from a first computing device to a second             |
| 4  | computing device across a network in order to continue execution on the second        |
| 5  | computing device, comprising:   |
| 6  | a receiving mechanism that is configured to receive the application                   |
| 7  | snapshot of the executing program from the first computing device on the second       |
| 8  | computing device, wherein the application snapshot includes a subprogram, an          |
| 9  | operand stack, and a point of execution;  |
| 10 | an examination mechanism that is configured to examine the application                |
| 11 | snapshot on the second computing device to identify the subprogram being              |
| 12 | executed and the point of execution within the subprogram wherein, the                |
| 13 | examination mechanism is configured to also examine the subprogram to                 |
| 14 | determine an expected structure of the operand stack at the point of execution;       |
| 15 | a validation mechanism that is configured to validate that the state of the           |
| 16 | application snapshot on the second computing device is consistent with the            |
| 17 | expected structure of the operand stack;  |
| 18 | a verifying mechanism configured to verify on the second computing                    |
| 19 | device that variables and argument arguments within the application snapshot are      |
| 20 | of the proper type; and   |
| 21 | an execution mechanism that is configured to resume execution of the                  |
| 22 | application snapshot on the second computing device at the point of execution         |
| 23 | from the first computing device if the state of the application snapshot is validated |
| 24 | as consistent with the expected structure of the operand stack.                       |

16. (Original) The apparatus of claim 15, wherein the examination mechanism includes a code verifier, wherein the code verifier is configured to ensure that:

| 4 | the subprogram does not cause the operand stack to overflow and                     |
|---|---|
| 5 | underflow;  |
| 6 | a use of a local variable does not violate type safety; and                         |
| 7 | an argument of an instruction is of an expected type.                               |
| 1 | 17. (Original) The apparatus of claim 15, wherein the operand stack                 |
| 2 | contains at least one local variable, at least one argument that is passed as a     |
| 3 | parameter to the subprogram, and an offset to the point of execution within the     |
| 4 | subprogram.   |
|   |   |
| 1 | 18. (Original) The apparatus of claim 16, wherein the expected structure of         |
| 2 | the operand stack includes a collective size of entries and the types of entries    |
| 3 | expected on the operand stack at the point of execution within the subprogram.      |
| 1 | 19 (Canceled).  |
| 1 | 20. (Original) The apparatus of claim 18, wherein the validation                    |
| 2 | mechanism is configured to ensure that the collective size of entries and the types |
| 3 | of entries on the operand stack agree with the collective size of entries and the   |
| 4 | types of entries expected on the operand stack.                                     |
| 1 | 21. (Original) The apparatus of claim 15, wherein in resuming execution             |
| 2 | of the application snapshot, the execution mechanism is configured to restart the   |
| 3 | subprogram at the point of execution within the second computing device.            |
| 1 | 22. (Previously presented) The method of claim 1, further comprising                |
| 2 | restoring the state of an object within the application snapshot on the second      |

- 3 computing device by changing a pointer from an address of the object on the first
- 4 computing device to an address of the object on the second computing device.
- 1 23. (Previously presented) The computer-readable storage medium of
- 2 claim 8, further comprising restoring the state of an object within the application
- 3 snapshot on the second computing device by changing a pointer from an address
- 4 of the object on the first computing device to an address of the object on the
- 5 second computing device.
- 1 24. (Previously presented) The apparatus of claim 15, further comprising
- 2 an object restoring mechanism that is configured to restore the state of an object
- 3 within the application snapshot on the second computing device by changing a
- 4 pointer from an address of the object on the first computing device to an address
- 5 of the object on the second computing device.